

Math 220C - Lecture 22

May 24, 2022

§ 1. Homological Methods & Sheaves (Part II)

Given $\alpha: \mathcal{F} \rightarrow \mathcal{G}$ morphism of sheaves, we define

i the sheaf $\text{Ker } \alpha$

ii the sheaf $\text{Coker } \alpha$

Kernel When $U \subseteq X$ open, set

$$\text{Ker } \alpha (U) = \text{Ker } \{ \alpha_U : \mathcal{F}(U) \rightarrow \mathcal{G}(U) \}$$

Restriction maps

$\text{Ker } \alpha (U) \rightarrow \text{Ker } \alpha (V)$ are naturally defined.

Check i $\text{Ker } \alpha$ is a sheaf

ii $(\text{Ker } \alpha)_* = \text{Ker } \{ \alpha_* : \mathcal{F}_* \rightarrow \mathcal{G}_* \}.$

Cokernel Presheaf

For $u \subseteq X$ open, define

$$\widetilde{\text{Coker } \alpha}(u) = \text{Coker} \{ \alpha_u : \mathcal{F}(u) \rightarrow \mathcal{G}(u) \}$$

Check \square $\widetilde{\text{Coker } \alpha}$ is a presheaf

$$\square \left(\widetilde{\text{Coker } \alpha} \right)_x = \text{Coker} (\alpha_x : \mathcal{F}_x \rightarrow \mathcal{G}_x)$$

Beware! $\widetilde{\text{Coker } \alpha}$ is not always a sheaf

Example $X = \mathbb{C}^*$, $0 \xrightarrow{\alpha} \mathbb{C}^*$, $\alpha(f) = e^{2\pi i f}$

Let $\widetilde{\mathcal{F}} = \widetilde{\text{Coker } \alpha}$

• $U = \mathbb{C} \setminus \mathbb{R}_{\leq 0}$, $\widetilde{\mathcal{F}}(U) = 0$

Indeed, U is *simply connected*, so logarithms

make sense $\Rightarrow \alpha$ is surjective in $U \Rightarrow \widetilde{\text{Coker } \alpha|_U} = 0$.

• $V = \mathbb{C} \setminus \mathbb{R}_{\geq 0} \Rightarrow \widetilde{\mathcal{F}}(V) = 0$

• $X = U \cup V = \mathbb{C}^*$. Note α is not surjective in X

since $\log 2$ is not defined in \mathbb{C}^* . $\Rightarrow \widetilde{\mathcal{F}}(X) \neq 0$.

Any $s \in \widetilde{\mathcal{F}}(X)$ restricts $s|_U = 0$, $s|_V = 0$ since

$\widetilde{\mathcal{F}}(U) = 0$, $\widetilde{\mathcal{F}}(V) = 0$. If $s \neq 0$ this contradicts

uniqueness of gluing $\Rightarrow \widetilde{\mathcal{F}}$ not a sheaf.

Sheafification

Goal

Given $\mathcal{F} \rightarrow X$ a presheaf, we define a sheaf $\mathcal{F}^\#$

& morphism of presheaves

$$i: \mathcal{F} \rightarrow \mathcal{F}^\#.$$

Remark In addition,

$$\boxed{\text{III}} \quad \mathcal{F} \text{ sheaf} \Rightarrow \mathcal{F} = \mathcal{F}^\#$$

$$\boxed{\text{IV}} \quad \mathcal{F}_* = \mathcal{F}_*^\# \quad \forall *$$

$\boxed{\text{V}}$ Given $\mathcal{F} \rightarrow \mathcal{G}$ we obtain $\mathcal{F}^\# \rightarrow \mathcal{G}^\#$ with

$$\begin{array}{ccc} \mathcal{F} & \longrightarrow & \mathcal{G} \\ \downarrow & & \downarrow \\ \mathcal{F}^\# & \longrightarrow & \mathcal{G}^\# \end{array}$$

commutative.

Definition

$\mathcal{F}^\#(U) = \left\{ (f_x)_{x \in U} \in \prod_{x \in U} \mathcal{F}_x, \text{ "locally compatible germs" i.e.} \right.$

$\left. \forall x \in U \exists x \in V \subseteq U, s \in \mathcal{F}(V) \text{ with } s_y = f_y \ \forall y \in V \right\}$

Example

\mathcal{F} = presheaf of constant functions

$\mathcal{F}^\#$ = sheaf of locally constant functions

Remark We define $\mathcal{F} \rightarrow \mathcal{F}^\#$ via

$$\mathcal{F}(U) \ni f \rightarrow (f_x)_x \in \mathcal{F}^\#(U).$$

Check $\mathcal{F}^\#$ is a sheaf & $\mathcal{F}_x = \mathcal{F}_x^\#$.

Conclusion — Cokernel sheaf

Given $\mathcal{F} \xrightarrow{\alpha} \mathcal{G}$, we define the cokernel sheaf:

(1) $\widetilde{\text{Coker } \alpha}$ presheaf

(2) sheafify $\text{Coker } \alpha := \widetilde{\text{Coker } \alpha}^\#$

Why does it work?

Assume $0 \rightarrow \mathcal{F} \xrightarrow{\alpha} \mathcal{G}$. We have by definition

$$\mathcal{G} \rightarrow \widetilde{\text{Coker } \alpha}$$

This gives

$$\mathcal{G} = \mathcal{G}^\# \rightarrow \widetilde{\text{Coker } \alpha}^\#$$

Then

$$0 \rightarrow \mathcal{F} \rightarrow \mathcal{G} \rightarrow \widetilde{\text{Coker } \alpha}^\# \rightarrow 0 \text{ exact}$$

as needed.

Exactness can be checked on stalks. We note that

$$\begin{aligned} \left(\widetilde{\text{Coker } \alpha}^\# \right)_x &= \left(\widetilde{\text{Coker } \alpha} \right)_x \\ &= \text{Coker}(\alpha_x : \mathcal{F}_x \rightarrow \mathcal{G}_x). \end{aligned}$$

$$\Rightarrow 0 \rightarrow \mathcal{F}_x \rightarrow \mathcal{G}_x \rightarrow \left(\widetilde{\text{Coker } \alpha}^\# \right)_x \rightarrow 0$$

exact, as needed.

§ 2. Flabby sheaves

\mathcal{F} is flabby provided $\forall v \subseteq u \subseteq X$ open,

$\mathcal{F}(u) \rightarrow \mathcal{F}(v)$ is surjective.

Remarks \square \mathcal{F} flabby $\Rightarrow \mathcal{F}/_u$ flabby $\forall u \subseteq X$ open

Indeed, for $v \subseteq w \subseteq u$,

$\mathcal{F}/_u(w) = \mathcal{F}(w) \rightarrow \mathcal{F}(v) = \mathcal{F}/_u(v)$ surjective.

\square Suffices to check $\forall u \subseteq X$ open

$\mathcal{F}(X) \rightarrow \mathcal{F}(u)$ surjective

Indeed, $\mathcal{F}(X) \rightarrow \mathcal{F}(u)$ shows $\mathcal{F}(v) \rightarrow \mathcal{F}(u)$
 $\swarrow \quad \nearrow$
 $\mathcal{F}(v)$ also surjective for $u \subseteq v$.

Key Lemma

If $0 \rightarrow \mathcal{F} \rightarrow \mathcal{G} \rightarrow \mathcal{H} \rightarrow 0$ exact, \mathcal{F} flabby then

\square $0 \rightarrow \mathcal{F}(u) \rightarrow \mathcal{G}(u) \rightarrow \mathcal{H}(u) \rightarrow 0$ exact, $\forall u \subseteq X$ open

\square \mathcal{F}, \mathcal{G} flabby $\Rightarrow \mathcal{H}$ flabby

Proof $\square \Rightarrow \square$

Let $v \subseteq u$ open. Compare the exact sequences

$$0 \rightarrow \mathcal{F}(u) \rightarrow \mathcal{G}(u) \rightarrow \mathcal{H}(u) \rightarrow 0$$

$$\downarrow \text{surj}$$

$$\downarrow \text{surj} \Rightarrow$$

$$\downarrow \text{surj}$$

$\Rightarrow \mathcal{H}$ flabby.

$$0 \rightarrow \mathcal{F}(v) \rightarrow \mathcal{G}(v) \rightarrow \mathcal{H}(v) \rightarrow 0$$