MA152 Solutions to Homework 4

May 13, 2017

1. We have

$$A^{-1} = \begin{pmatrix} 1/2 & 0 & 0 & 0\\ 0 & 1/3 & 0 & 0\\ 0 & 0 & 1 & 0\\ 0 & 0 & 0 & 1/3 \end{pmatrix}$$

So $\mathbf{1}^T A^{-1} \mathbf{1} = (1/2 + 1/3 + 1 + 1/3) = 13/6$. This is not zero, so we can use the nonsingular matrix theorem. The value of the game is $V = (\mathbf{1}^T A^{-1} \mathbf{1})^{-1} = 6/13$. The optimal strategies for each players are given by

 $\mathbf{p}^T = V \mathbf{1}^T A^{-1} = \begin{pmatrix} 3/13 & 2/13 & 6/13 & 2/13 \end{pmatrix}$

and

 $\mathbf{q}^{T} = V \mathbf{1} A^{-1} = \begin{pmatrix} 3/13 & 2/13 & 6/13 & 2/13 \end{pmatrix}$

2. This is an upper triangular matrix, so let us try looking for an equalizing strategy. Let Player I's strategy be $\mathbf{p}^T = (p_1 \ p_2 \ p_3 \ p_4)$ and let Player II's strategy be $\mathbf{p}^T = (q_1 \ q_2 \ q_3 \ q_4)$. We get the following system of equations for Player I:

$$p_{1} = V$$

$$-p_{1} + p_{2} = V$$

$$-2p_{2} + p_{3} = V$$

$$-p_{1} + p_{2} - p_{3} + p_{4} = V$$

$$p_{1} + p_{2} + p_{3} + p_{4} = 1$$

Solving this, we get V = 1/13 and

$$\mathbf{p}^T = \begin{pmatrix} 1/13 & 2/13 & 5/13 & 5/13 \end{pmatrix}$$

For Player II, we get the system

$$\begin{array}{ccc} q_1 - q_2 & -q_4 = V \\ q_2 - 2q_3 + q_4 = V \\ q_3 - q_4 = V \\ q_4 = V \\ q_1 + q_2 + q_3 + q_4 = 1 \end{array}$$

Solving this system, we get

$$\mathbf{q}^T = \begin{pmatrix} 6/13 & 4/13 & 2/13 & 1/13 \end{pmatrix}$$

3. Denote the columns by A,B and C, and denote the rows by 1, 2 and 3. Now, swapping rows 1 and 3 and columns A and C leaves the payoff matrix unchanged. So, by Invariance, we can assume that the optimal strategies **p** and **q** satisfy $\mathbf{p}(1) = \mathbf{p}(3)$ and $\mathbf{q}(A) = \mathbf{q}(C)$.

We look for an equalizing strategy. For Player I, we get

$$-4\mathbf{p}(1) + \mathbf{p}(2) + 2\mathbf{p}(3) = V$$

$$\mathbf{p}(1) - 5\mathbf{p}(2) + \mathbf{p}(3) = V$$

$$2\mathbf{p}(1) + \mathbf{p}(2) - 4\mathbf{p}(3) = V$$

$$\mathbf{p}(1) + \mathbf{p}(2) + \mathbf{p}(3) = 1$$

Now applying $\mathbf{p}(1) = \mathbf{p}(3)$, we get

$$-2\mathbf{p}(1) + \mathbf{p}(2) = V$$

$$2\mathbf{p}(1) - 5\mathbf{p}(2) = V$$

$$2\mathbf{p}(1) + \mathbf{p}(2) = 1$$

Solving this system, wet get V = -1/2 and

$$\mathbf{p}^T = \begin{pmatrix} 3/8 & 1/4 & 3/8 \end{pmatrix}$$

Since the matrix is symmetric, the system of equations for finding an equalizing strategy for Player II is the same as for Player I. So we also have

$$\mathbf{q}^T = \begin{pmatrix} 3/8 & 1/4 & 3/8 \end{pmatrix}$$

4. Answer 1: Note that the payoff matrix A in this question satisfies A(i, j) = i - j. So we have

$$A(i, j) = i - j = -(j - i) = -A(j, i)$$

and so A is skew-symmetric. Therefore the value of the game is zero.

Answer 2: We can note that the bottom right entry of the matrix (entry (1000, 1000)) is a saddle point, since A(1000, 1000) = 0, and 0 is the smallest number in the last row and the largest number in the last column. So V = 0.

Answer 3: Note that the last row dominates all other rows. Similarly, the last column dominates all other columns. All that remains is the bottom right entry, which is 0. So V = 0.

5. (a) Given this strategy for Player II, the possible payoffs are

$$A\mathbf{q} = \begin{pmatrix} 2q+1\\ -4q+5 \end{pmatrix}$$

Player I will choose the first row if $2q + 1 > -4q + 5 \Leftrightarrow q > 2/3$, and the second row otherwise.

(b) Player II knows that this is what Player I will choose. So Player II will choose the value of q that maximizes the outcome.

If $0 \le q \le 2/3$, then the payoff is -4q + 5. Along this interval, this function is minimized at q = 2/3. The payoff in this case is 7/3.

If $2/3 < q \leq 1$, then the payoff is 2q + 1, which is always greater than 7/3 on this interval. So the minimum payoff that Player II can achieve is 7/3, if they choose q = 2/3.

6. (a) Let C denote the matrix whose entries are all c. So B = A + C.
Answer 1: Let p, q be strategies chosen by players I and II during a game of B. Then the payoff is

$$\mathbf{p}^T B \mathbf{q} = \mathbf{p}^T (A + C) \mathbf{q} = \mathbf{p}^T A \mathbf{q} + c$$

where the last equality follows from the fact that the entries of \mathbf{p} and \mathbf{q} sum to 1.

In particular, playing the game B is exactly the same as playing the game A, and then adding c to the payoff of A. It is clear that the optimal strategy for both players in such a game is just to use their optimal strategies for A. In this case the payoff is V + c, and so this is the value of B.

Answer 2: Let \mathbf{p}, \mathbf{q} be a pair of optimal strategies for A. Then, from the definition, we get that

$$\min(\mathbf{p}^T A) = V = \max(A\mathbf{q})$$

(where min and max here represent the min and max entry of a vector).

Now we have

$$\min(\mathbf{p}^T B) = \min(\mathbf{p}^T (A + C)) = \min(p^T A + \mathbf{p}^T C)$$

Note that since \mathbf{p} is a probability vector, every entry of the vector $\mathbf{p}^T C$ is just c. So

$$\min(\mathbf{p}^T B) = \min(\mathbf{p}^T A) + c = V + c$$

Similarly,

$$\max(B\mathbf{q}) = V + c$$

From the definition of value, we get that V + c is the value of B, with optimal strategies **p** and **q**.

(b) Let A be the matrix given in the question. Notice that

The first matrix on the right hand side is skew-symmetric, so has value 0. From (a), we get that

$$\operatorname{Val}(A) = 0 + 2 = 2$$