

Then (Lagrange inversion formula). Let  $G(x)$  be formal power series s.t.  $G(0) \neq 0$ . Let  $A(x)$  formal power series s.t.  $A(0) = 0$ , and

$$A(x) = x G(A(x)).$$

Then for  $k, n$  we have

$$n[x^n] A(x)^k = k [x^{n-k}] (G(x))^n.$$

Ex.  $C(x) = \sum_{n \geq 0} C_n x^n$  (Catalan)

satisfies  $C(x) = 1 + x C(x)^2$ . Define  $A(x) = C(x) - 1$

$$A(x) \neq 1 + x(A(x) + 1)^2$$

$$A(x) = x(A(x) + 1)^2, \quad G(x) = (1+x)^2$$

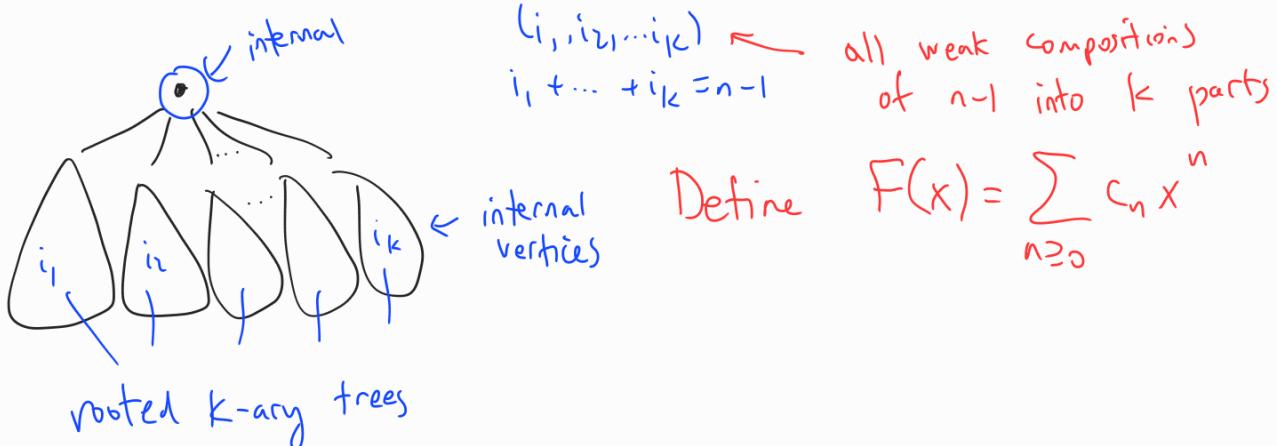
$$n[x^n] A(x) = [x^{n-1}] (1+x)^{2n} = \binom{2n}{n-1}$$

$$\text{If } n > 0, \quad C_n = \frac{1}{n} \binom{2n}{n-1} = \frac{(2n)!}{n \cdot (n-1)! \cdot (n+1)!} = \frac{1}{n+1} \binom{2n}{n}$$

Ex.  $C_n = \# \text{ rooted binary trees w/ } (n+1) \text{-leaves}$   
 $= \# \text{ rooted binary trees w/ } n \text{ internal vertices.}$

Let  $C_n = \# \text{ rooted k-ary trees w/ } n \text{ internal vertices}$

Recurrence relation:  $C_n = \sum_{\substack{(i_1, i_2, \dots, i_k) \\ i_1 + \dots + i_k = n-1}} c_{i_1} c_{i_2} \dots c_{i_k} \quad \text{if } n > 0$



Define  $F(x) = \sum_{n \geq 0} C_n x^n$

$$F(x) = 1 + \sum_{n>0} c_n x^n = 1 + x \sum_{n>0} \sum_{\substack{i_1, \dots, i_k \\ i_1 + \dots + i_k = n-1}} c_{i_1, \dots, i_k} x^{n-1} = 1 + x F(x)^k$$

$$\text{Define } A(x) = F(x) - 1 : \quad A(x) + 1 = 1 + x(A(x) + 1)^k$$

$$G(x) = (1+x)^k$$

$$\text{Lagrange: } [x^n] A(x) = [x^{n-1}] (1+x)^{kn} = \binom{kn}{n-1}$$

$$c_n = \frac{1}{n} \binom{kn}{n-1} = \frac{1}{(k-1)n+1} \binom{kn}{n}.$$

### Proof of Lagrange inversion

Lemma. If  $F(0) = 0$ , then  $\exists G(x)$  s.t.  $G(F(x)) = x$

$$\Leftrightarrow [x^1] F(x) \neq 0.$$

Formal Laurent series: Series of the form  $\sum_{n \geq N} a_n x^n$  where  $N$  is some integer (depending on the series)

Only finitely many negative powers at a given time.

Addition / multiplication defined as before

Every  $\underset{n \neq 0}{\text{nonzero}}$  formal Laurent series has a multiplicative inverse.

$$\text{derivative } D \sum_{n \geq N} a_n x^n = \sum_{n \geq N} n a_n x^{n-1}$$

$$\text{Note } [x^{-1}] D F = 0$$

$\forall$  formal Laurent series

$$\text{Start w/ } A(x) = x G(A(x)). \quad (*)$$

$$[x^1] A(x) = [x^0] G(A(x)) = G(0) \neq 0$$

$$\Rightarrow \exists B(x) \text{ s.t. } B(A(x)) = x = A(B(x)).$$

$$\text{Substitute } x \rightarrow B(x) \text{ in } (*); \quad x = B(x) G(x).$$

Write  $A(x)^k = \sum_{d \geq k} \alpha_d x^d$ . Substitute  $x \rightarrow B(x)$  to get

$$x^k = \sum_{d \geq k} \alpha_d B(x)^d \quad \text{Take derivative}$$

$$k x^{k-1} = \sum_{d \geq k} \alpha_d d \cdot B(x)^{d-1} B'(x) \quad \begin{array}{l} \text{Divide by } B(x)^n \\ (\text{use formal Laurent series}) \end{array}$$

$$\frac{k x^{k-1}}{B(x)^n} = \sum_{d \geq k} d \alpha_d B(x)^{d-n-1} B'(x) \quad \text{Take } [x^{-1}]$$

If  $d \neq n$ , then  $\underbrace{B(x)^{d-n-1} B'(x)}_{\text{no } [x^{-1}] \text{ term}} = \left( \frac{B(x)^{d-n}}{d-n} \right)'$

$$[x^{-1}] \frac{k x^{k-1}}{B(x)^n} = [x^{-1}] n \alpha_n \frac{B'(x)}{B(x)} \stackrel{\text{exercise}}{=} n \alpha_n = n [x^n] A(x)^k$$

$$\Rightarrow [x^{-1}] \frac{k x^{k-1}}{B(x)^n} = [x^{-1}] k x^{k-n-1} G(x)^n = k [x^{k-n}] G(x)^n. \quad \square$$