

What is this course about?

3 sample problems

① Linear recurrence relations:

Fibonacci sequence: $0, 1, 1, 2, 3, 5, 8, 13, \dots$

Can I find formula (non-recursive)?

② Counting # ways to write n pairs of balanced parentheses

$n=2$: $()()$ $(())$

$n=3$: $()()()$ $(())()$ $()(())$ $((()))$ $(())()$

general n ?

③ How many permutations of n things have no fixed points?

Bijections Given sets X, Y , functions $f: X \rightarrow Y$
 $g: Y \rightarrow X$

we say inverses of each other if: $f \circ g = id_Y$ & $g \circ f = id_X$

i.e., For all $y \in Y$, $f(g(y)) = y$, For all $x \in X$, $g(f(x)) = x$

If so, f and g are bijections.

Prop. If a bijection exists between X, Y then $|X| = |Y|$.

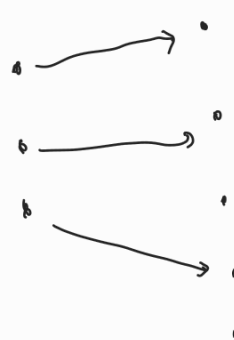
If f is bijection, then it is

- one-to-one (injective), i.e., if $f(x) = f(x')$, then $x = x'$
- onto (surjective), i.e., for all $y \in Y$, there is some x so that $f(x) = y$

f is bijection \iff injective & surjective.

f injective $\Rightarrow |X| \leq |Y|$

f surjective $\Rightarrow |X| \geq |Y|$

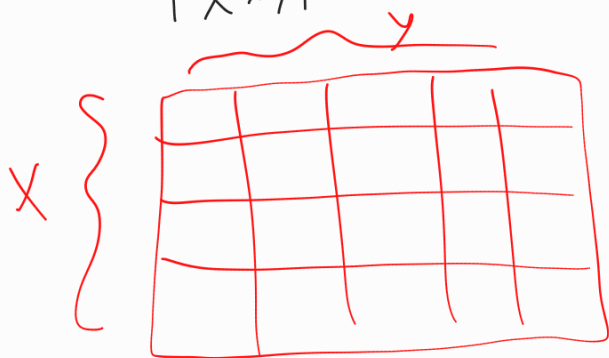


Sum principle: X, Y sets w/ no overlap.

Then $|X \cup Y| = |X| + |Y|$.

Product principle: $X \times Y = \{(x, y) \mid x \in X, y \in Y\}$

$$|X \times Y| = |X| \cdot |Y|$$



12-fold way: Assigning k balls to n boxes.

subject to some condition:

Think of assignment as function $f: \{\text{balls}\} \rightarrow \{\text{boxes}\}$

Conditions on f :

- ① f is injective
- ② f is surjective
- ③ no condition

Conditions on balls:

- ① all identical
- ② considered different

Conditions on boxes:

- ① all identical - indistinguishable
- ② considered different - distinguishable

balls / boxes	f arbitrary	f injective	f surjective
dist / dist			
in dist / dist			
dist / in dist		$\begin{cases} 1 & \text{if } n \geq k \\ 0 & \text{if } n < k \end{cases}$	
in dist / in dist		$\begin{cases} 1 & \text{if } n \geq k \\ 0 & \text{if } n < k \end{cases}$	

Weak

Induction : Prove sequence of statements $P(0), P(1), P(2), \dots$

Strategy: ① Prove $P(0)$ is true. Base case
 ② Use $P(n)$ to prove $P(n+1)$ is true. Inductive step

Eg. $P(n)$ is statement " $\sum_{i=0}^n i = \frac{n(n+1)}{2}$ "

① $P(0)$: " $\sum_{i=0}^0 i = \frac{0 \cdot (1)}{2} = 0$ " ✓

② $\sum_{i=0}^{n+1} i = \sum_{i=0}^n i + (n+1) = \frac{n(n+1)}{2} + (n+1)$
 $= (n+1) \left(\frac{n}{2} + 1 \right) = (n+1) \left(\frac{n+2}{2} \right)$

$P(n) \Rightarrow P(n+1)$ ✓

$\left(\frac{n}{2} + 1 \right) = \frac{n}{2} + \frac{2}{2} = \frac{n+2}{2}$
 $\rightarrow \frac{(n+1)(n+2)}{2}$

Subsets. S set, T is a subset if every element of T belongs to S .

Note: T could be empty, and also $T=S$ ✓.

Thm. If $|S|=n$, there are 2^n subsets of S .

pf. Base case: $n=0$, $S=\emptyset$, only subset is $T=\emptyset$
 $2^0=1$ ✓

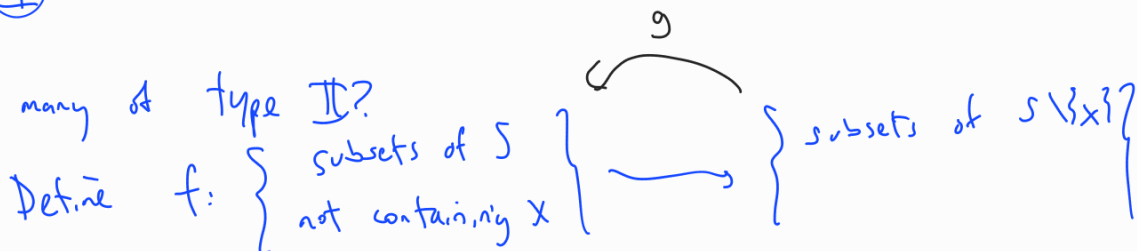
Inductive step: Suppose $|S|=n+1 > 0$.

Pick some element $x \in S$. Consider 2 kinds of subsets

Ⓘ: T contains x

Ⓜ: T does not contain x

How many of type Ⓜ?



$f(T) = T$, $g(T) = T$, get bijection

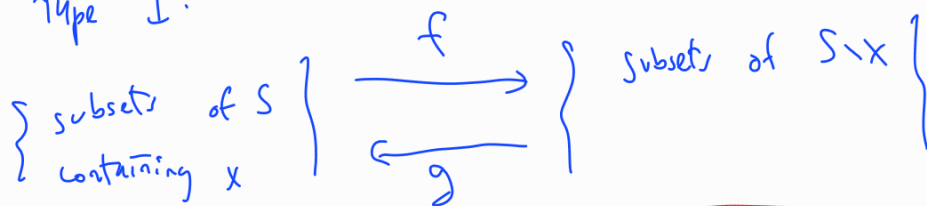
$\Rightarrow 2^n$ subsets of type Ⓜ.

EX. $S = \{a, b, c\}$, $x = b$

subsets of S not containing $b = \left. \begin{array}{l} \emptyset \\ \{a\} \\ \{a, c\} \\ \{c\} \end{array} \right\}$

subsets of $\{a, c\} = \left. \begin{array}{l} \emptyset \\ \{a\} \\ \{a, c\} \\ \{c\} \end{array} \right\}$

How many of type Ⓘ?



$f(T) = T \setminus x$, $g(U) = U \cup \{x\}$

bijection $\Rightarrow 2^n$ type Ⓘ subsets.

subsets = # type Ⓘ + # type Ⓜ = $2^n + 2^n = 2(2^n) = 2^{n+1}$ □

T subset of S not
containing x

$$g(f(T)) = g(T \setminus x) = (T \setminus x) \cup \{x\}$$

U subset of $S \setminus x$:

$$f(g(U)) = f(U \cup \{x\}) = U$$

Ex. $\sum_{i=0}^n i^2 = \frac{n(n+1)(2n+1)}{6}$

$\sum_{i=0}^n i^3 = ??$ deg 4 polynomial
in n

$\sum_{i=0}^n i^d =$ deg $d+1$ polynomial in n .

$S = \{a, b, c\}, x = b$

~~$\{b\}$~~

~~$\{a, b\}$~~

~~$\{b, c\}$~~

~~$\{a, b, c\}$~~

\emptyset

$\{a\}$

$\{c\}$

$\{a, c\}$