Math 475, Fall 2015
Homework 10
Due: Friday, Dec. 4
Notation: $[n]=\{1, \ldots, n\}$.
Recall from class that the girth of a graph is the shortest length of a cycle it contains, and that we proved that if $G$ is a simple planar graph with $n$ vertices, $m$ edges, and finite girth $g$, then

$$
\begin{equation*}
m \leq \frac{g}{g-2}(n-2) \tag{}
\end{equation*}
$$

(1) In this exercise, you'll find some examples where $\left(^{*}\right)$ is an equality.
(a) Draw an example of a simple planar graph of girth 3 with 7 vertices and 15 edges.
(b) For $n \geq 4$, show that the complete bipartite graph $K_{2, n-2}$ is planar, has girth 4, and $\left(^{*}\right)$ is an equality.
(c) Draw an example of a simple planar graph of girth 5 with 8 vertices and 10 edges.
(2) The following graph $G$ is the Petersen graph:

(a) Find 5 edges to contract so that the resulting graph is isomorphic to the complete graph $K_{5}$. Conclude that $G$ is not planar.
(b) Compute the girth of $G$ and show that the inequality $\left({ }^{*}\right)$ fails. (This gives a different reason for the fact that $G$ is not planar.)
(3) If $r \leq n$, an $r \times n$ Latin rectangle is an $r \times n$ matrix with entries in [ $n$ ] such that no row and no column has a repeated entry. In this exercise, you will prove the following theorem of Ryser: if $r<n$, then any $r \times n$ Latin rectangle can be extended to an $(r+1) \times n$ Latin rectangle (meaning that you can add an extra row at the bottom with the numbers $1, \ldots, n$ so that the result is an $(r+1) \times n$ Latin rectangle).
(a) Let $G=(X, Y)$ be a simple bipartite graph such that $|X| \leq|Y|$. Assume that there is a positive integer $d$ such that $\operatorname{deg}(v)=d$ for all vertices $v$ of $G$. Show that there is a perfect matching of $X$ into $Y$.
(b) Given an $r \times n$ Latin rectangle, build a bipartite graph $G$ on $2 n$ vertices with $X=Y=[n]$ where $i \in X$ is connected to $j \in Y$ if and only if $i$ does not appear in the $j$ th column of the $r \times n$ Latin rectangle. Show that every vertex of $G$ has degree $n-r$ and use (a) to conclude that $G$ has a perfect matching. Use the perfect matching to construct the next row.
(4) Let $G$ be a simple planar graph with $n$ vertices and $m$ edges.
(a) Assume that $\operatorname{deg}(v) \geq 5$ for all vertices $v$ (this implies $n \geq 6$ ). Use the bound $m \leq 3 n-6$ to prove that $n \geq 12$.
(b) Use (a) to show that if $m<30$, then $G$ has a vertex of degree $\leq 4$.
(c) Use (b) and modify the proof of the five-color theorem to show that if $m<30$, then $G$ has a proper 4-coloring. Don't cite the four-color theorem! (So this proves the four-color theorem for planar graphs with $<30$ edges.)

Remark. You can't improve the bounds in $\# 4(\mathrm{a}, \mathrm{b})$ since there is a planar graph with 12 vertices such that every vertex has degree 5 (and thus has 30 edges). To see one, look here: https://upload.wikimedia.org/wikipedia/commons/8/83/Icosahedron_graph.svg

